

METHOD OF CONTROLLING THE FLOW OF INFORMATION  
BETWEEN SENDERS AND RECEIVERS ACROSS  
LINKS BEING USED AS CHANNELS

CROSS-REFERENCE TO RELATED APPLICATIONS

5        This application contains subject matter which  
is related to the subject matter of the following  
applications, each of which is assigned to the same  
assignee as this application and filed on the same  
day as this application. Each of the below listed  
10 applications is hereby incorporated herein by  
reference in its entirety:

      "SYSTEM OF CONTROLLING THE FLOW OF INFORMATION  
BETWEEN SENDERS AND RECEIVERS ACROSS LINKS BEING USED  
AS CHANNELS," by Gregg et al., Serial No. \_\_\_\_\_,  
15 (Docket No. PO9-98-124); and

      "CONTROLLING THE FLOW OF INFORMATION BETWEEN  
SENDERS AND RECEIVERS ACROSS LINKS BEING USED AS  
CHANNELS," by Gregg et al., Serial No. \_\_\_\_\_,  
(Docket No. PO9-98-125).

20

TECHNICAL FIELD

      This invention relates, in general, to data  
communication and, in particular, to controlling the  
flow of information between senders of information  
and receivers of that information, when the  
25 information is being forwarded across links being  
used as channels.

PO9-98-109-1

BACKGROUND ART

Communication between computer systems usually involves a sending system (sender) sending a command to a receiving system (receiver) over a link used to couple the sender and the receiver. The receiving system then, typically, sends a response back over the link to the sending system.

One example of a link used by International Business Machines Corporation to communicate between a sender and a receiver is an intersystem channel link. In particular, an intersystem channel link couples an intersystem channel adapter on one system (e.g., a central processor) that may be used for sending or receiving messages with an intersystem channel adapter on another system (e.g., a coupling facility that contains data shared by any central processor coupled thereto) that may also be used for sending or receiving messages.

An intersystem channel link supports connections up to approximately 20 kilometers, and the propagation time on the link (at 20 kilometers), when using fiber optic conductors, is approximately 100 microseconds in each direction. To minimize message latency, a minimum number of round trip acknowledgements is used, and thus, large data buffers at each end of the link are needed for flow control. In one example, a message exchange requires only one single round trip over the link.

As more and more data is desired to be included with the message, larger data buffers are needed to

FOIA b 7 - D

maintain the single round trip flow control.  
However, eventually, the expense of large buffers  
becomes prohibitive and the flow control is modified  
to add intermediate acknowledgement exchanges to  
5 throttle the data transfer preventing the buffers  
from overrunning. These extra exchanges, though,  
significantly increase message latency.

Based on the foregoing, a need exists for a more  
direct communication link between senders and  
10 receivers of information. Further, a need exists for  
a communication link that does not need to couple  
channel adapters, but still may act as a channel.  
That is, a need exists for a communication protocol  
in which the advantages of a channel (e.g., having a  
15 direct memory adapter engine and offering protection  
of memory) may be realized. Additionally, a need  
exists for a communication protocol that does not  
require large data buffers and does not significantly  
add to message latency.

20 SUMMARY OF THE INVENTION

The shortcomings of the prior art are overcome  
and additional advantages are provided through the  
provision of a method of controlling the flow of  
information across links between senders and  
25 receivers. The method includes, for instance,  
sending a request from a sender to a receiver over a  
link; forwarding a data request indication from the  
receiver to the sender in response to the request, in  
which the data request indication indicates that the  
30 receiver is prepared to receive further information  
from the sender; and providing, by the sender, in

response to the data request indication, the further information. While the receiver is preparing for the further information, the link is being used as a channel in that the link is not stopped.

5 In another aspect of the present invention, a method of controlling the flow of information across links between senders and receivers is provided. The method includes, for instance, including in a packet a sequence number usable in maintaining delivery  
10 order of the packet. The packet, however, does not include a memory address and does not require an explicit individual response. The packet is sent from a sender to a receiver across a link, and the sequence number is used to determine if the packet is  
15 in proper order for processing by the receiver.

Advantageously, the flow control capabilities of the present invention provide a more direct communication path between senders and receivers. Further, the flow control protocol does not require  
20 large data buffers and offers low latency messaging. Thus, the present invention is less expensive than other communication protocols.

Additional features and advantages are realized through the techniques of the present invention.  
25 Other embodiments and aspects of the invention are described in detail herein and are considered a part of the claimed invention.

Variable	Mean	SD	Min	Max
Age	38.5	12.5	18	65
Gender	Male	Female		
Marital status	Married	Single		
Education	High school	College		
Occupation	Manager	Worker		
Income	\$20,000	\$40,000		
Health status	Good	Poor		
Exercise frequency	Weekly	Monthly		
Stress level	Low	High		
Sleep quality	Good	Poor		
Dietary habits	Healthy	Unhealthy		
Alcohol consumption	None	Occasional		
Tobacco use	Non-smoker	Smoker		
Family size	2	3		
Work hours	40	50		
Commuting time	30	45		
Home ownership	Owner	Renter		
Neighborhood safety	Safe	Unsafe		
Access to green spaces	Yes	No		
Proximity to schools	Close	Far		
Public transportation	Good	Poor		
Crime rate	Low	High		
Weather conditions	Good	Poor		
Local economy	Strong	Weak		
Community involvement	High	Low		
Local government	Effective	Ineffective		
Local culture	Diverse	Homogeneous		
Local history	Rich	Poor		
Local infrastructure	Good	Poor		
Local services	Good	Poor		
Local amenities	Good	Poor		
Local environment	Good	Poor		
Local safety	Good	Poor		
Local quality of life	Good	Poor		
Local satisfaction	High	Low		
Local happiness	High	Low		
Local well-being	High	Low		
Local health	Good	Poor		
Local wealth	High	Low		
Local power	High	Low		
Local influence	High	Low		
Local status	High	Low		
Local reputation	High	Low		
Local image	Good	Poor		
Local brand	Strong	Weak		
Local identity	Strong	Weak		
Local loyalty	High	Low		
Local commitment	High	Low		
Local engagement	High	Low		
Local participation	High	Low		
Local involvement	High	Low		
Local contribution	High	Low		
Local impact	High	Low		
Local legacy	High	Low		
Local heritage	High	Low		
Local tradition	High	Low		
Local custom	High	Low		
Local habit	High	Low		
Local practice	High	Low		
Local routine	High	Low		
Local habit	High	Low		
Local behavior	High	Low		
Local action	High	Low		
Local response	High	Low		
Local reaction	High	Low		
Local effect	High	Low		
Local result	High	Low		
Local outcome	High	Low		
Local consequence	High	Low		
Local impact	High	Low		
Local influence	High	Low		
Local power	High	Low		
Local authority	High	Low		
Local control	High	Low		
Local management	High	Low		
Local organization	High	Low		
Local structure	High	Low		
Local system	High	Low		
Local framework	High	Low		
Local foundation	High	Low		
Local base	High	Low		
Local core	High	Low		
Local center	High	Low		
Local hub	High	Low		
Local node	High	Low		
Local point	High	Low		
Local spot	High	Low		
Local site	High	Low		
Local place	High	Low		
Local area	High	Low		
Local region	High	Low		
Local territory	High	Low		
Local domain	High	Low		
Local jurisdiction	High	Low		
Local sphere	High	Low		
Local range	High	Low		
Local scope	High	Low		
Local scale	High	Low		
Local level	High	Low		
Local grade	High	Low		
Local rank	High	Low		
Local position	High	Low		
Local status	High			

5

10

bus adapter of FIG. 1;

15

20

the present invention;

25

FIG. 6 depicts one example of a self-timed interface header for channel adapter memory responses, in accordance with the principles of the present invention;

5           FIG. 7 depicts one embodiment of the flow of packets for an I/O input operation, in accordance with the principles of the present invention;

10           FIG. 8 depicts one embodiment of the flow of packets for an I/O output operation, in accordance with the principles of the present invention;

FIG. 9a depicts one embodiment of a frame sent across an intersystem channel;

15           FIG. 9b depicts one example of a header of the frame of FIG. 9a;

20           FIGS. 10a-10c illustrate examples of the flow of messages between an OS/390 processor and a coupling facility, in accordance with the principles of the present invention;

25           FIGS. 11a-11b illustrate additional examples of the flow of messages between an OS/390 processor and a coupling facility, in accordance with the principles of the present invention;

PO9-98-109

FIG. 12 depicts an example of the components of a memory bus adapter of the present invention;

5 FIG. 13 depicts one example of a header used for an integrated cluster bus (ICB) flow control, in accordance with the principles of the present invention;

10 FIG. 14 depicts an example of the components of the ICB control logic depicted in FIG. 12, in accordance with the principles of the present invention;

15 FIG. 15 depicts one embodiment of the logic associated with a write data example in which a message is sent from an OS/390 processor to a coupling facility using an integrated cluster bus, in accordance with the principles of the present invention;

20 FIG. 16 depicts one embodiment of the logic associated with a read data example in which a message is sent from an OS/390 processor to a coupling facility using an integrated cluster bus, in accordance with the principles of the present invention;

25 FIG. 17 depicts one embodiment of the logic associated with error checking performed when a command is received by a processor, in accordance with the principles of the present invention;

1050990" 8844/860

FIG. 18 depicts one embodiment of the logic associated with error checking performed when a packet is received over an ICB self-timed interface link, in accordance with the principles of the present invention;

FIG. 19 depicts one embodiment of the logic associated with sequence error checking performed when a packet is received over an ICB self-timed interface link, in accordance with the principles of the present invention;

FIG. 20 depicts one embodiment of the logic associated with checking the order of commands received from processors, in accordance with the principles of the present invention; and

FIG. 21 depicts one embodiment of the logic associated with checking the order of packets received over a link, in accordance with the principles of the present invention.

BEST MODE FOR CARRYING OUT THE INVENTION

In accordance with the principles of the present invention, the flow of information between senders and receivers is being controlled across links that are being used as channels.

One embodiment of a computing environment incorporating and using the flow control capabilities of the present invention is depicted in FIG. 1 and described in detail herein. Computing environment 100 is based, for instance, on the Enterprise Systems



Architecture (ESA)/390 offered by International Business Machines Corporation. ESA/390 is described in an IBM Publication entitled Enterprise Systems Architecture/390 Principles of Operation, IBM

- 5 Publication No. SA22-7201-04, June 1997, which is hereby incorporated herein by reference in its entirety.

- 10 Computing environment 100 includes, for instance, one or more central processing complexes (CPCs) 102 coupled to at least one coupling facility 104, each of which is described below.

- 15 Each central processing complex 102 includes one or more central processors 106 (central processing units) and main memory 108, which is accessed by the central processors. Each central processor may execute an operating system, such as the OS/390 or Multiple Virtual Storage (MVS)/ESA operating system offered by International Business Machines Corporation.

- 20 Coupling facility 104 is a sharable facility, which includes one or more central processors 110 and storage 112, which is indirectly accessible by processors 106 and directly accessible by processors 110. Each central processor 110 may execute coupling facility control code to perform operations requested by processors 106. In one embodiment, coupling facility 104 is a structured-external storage processor (SES).

- 30 Aspects of the operation of a coupling facility are described in detail in such references as Elko et

al., U.S. Patent No. 5,317,739, entitled "Method and Apparatus For Coupling Data Processing Systems", issued May 31, 1994; Elko et al., U.S. Patent No. 5,561,809, entitled "In A Multiprocessing System Having A Coupling Facility, Communicating Messages Between The Processors And The Coupling Facility In Either A Synchronous Operation Or An Asynchronous Operation", issued on October 01, 1996; Elko et al., U.S. Patent No. 5,706,432, entitled "Mechanism For Receiving Messages At A Coupling Facility", issued January 06, 1998; and the patents and applications referred to therein, all of which are hereby incorporated herein by reference in their entirety.

Multiple central processing complexes and coupling facilities may be interconnected. Further, each central processing complex and/or coupling facility processors may be logically partitioned, such that it can run multiple instances of the OS/390 and/or coupling facility control code.

Central processing complex 102 and coupling facility 104 are each coupled to a memory bus adapter 114. Memory bus adapter 114 provides assistance in fetching data from and storing data in memory. In particular, the memory bus adapter is used to couple the memory to various peripheral components, such as input/output (I/O) devices 116. As examples, the input/output devices include disk drives, tape drives, local area network (LAN) attachments, and wide area network (WAN) attachments.

In one embodiment, in order to couple the memory bus adapters to the I/O devices, each memory bus

adapter 114 is coupled to one or more channel  
adapters 118, 120 via memory interfaces, referred to  
as self-timed interfaces (STIs) 124. The channel  
adapters then provide the I/O interface attachment  
5 points to the I/O devices. As examples, one of the  
channel adapters is a parallel channel adapter 118,  
which provides a parallel channel interface 126 to  
I/O devices 116; and another channel adapter is an  
ESCON channel adapter 120, which provides an ESCON  
10 channel interface 128 to I/O devices 116. In  
addition to the above, an intersystem channel (ISC)  
adapter 122, which provides an intersystem channel  
interface 130 to, for example, another central  
processing complex, is also coupled to the memory bus  
15 adapter via STI 124.

Channel adapters 118, 120, 122 are, for  
instance, direct memory adapter (DMA) engines, which  
move data to and from memory 108, 112 without any  
direct involvement of any of processors 106, 110.  
20 They also provide protection from the peripherals or  
other components storing in any arbitrary location in  
memory. The peripherals or other components can only  
store into the memory locations known to the channel  
adapters, and those locations are determined by the  
25 programs running in processors 106, 110. The channel  
adapters are considered a trusted, integral part of  
the central processing complex.

As described above, the channel adapters are  
coupled to memory bus adapter 114 via self-timed  
30 interface 124. In one example, self-timed interface  
124 includes a link that operates at relatively short  
distances (about 10 meters). The link is comprised

09874438-060504

of eight data signals, one combination parity and tag  
signal, and a clock signal in each direction. The  
maximum link propagation delay is about 40  
nanoseconds, and thus, the time it takes to transmit  
5 a single 128 byte memory line is about 500  
nanoseconds (including the packet overhead of headers  
and checking fields).

Coupled to one end of the STI link is memory bus  
adapter 114, which is described in further detail  
10 with reference to FIG. 2. Memory bus adapter 114  
includes one or more STI macros 200 and switch logic  
202. In one embodiment, each macro is coupled to one  
end of a self-timed interface link 124.

Each STI macro is responsible for retiming the  
15 data and parity/tag signals of the link to the clock  
signal. The retiming compensates for propagation  
time variations of the various signals so that the  
data and parity/tag signals can be coherently sampled  
by the clock signal. The propagation time variations  
20 are due to differences in the signal drivers,  
receivers and wire. (STI is further described in  
U.S. Patent 5,757,297, entitled "Method and Apparatus  
for Recovering a Serial Data Stream Using a Local  
Clock", Ferraiolo et al., issued on May 26, 1998; and  
25 co-pending U.S. Patent Application No. 08/660,648,  
entitled "Adaptive Filtering Method and Apparatus to  
Compensate For a Frequency Difference Between Two  
Clock Sources", Ferraiolo et al., filed on June 7,  
1996, each of which is assigned to International  
30 Business Machines Corp. and each of which is hereby  
incorporated herein by reference in its entirety).

005050" 88442860

STI macros 200 are also responsible for transmitting, receiving, buffering, and regulating packets and control sequences over STI. Examples of packets managed by STI are depicted in FIGs. 3a-3b and described below.

Packet 300a (FIG. 3a) includes, for instance, a header 302a and a longitudinal redundancy check 304a. Header 302a has a variable length (e.g., 8 or 16 bytes) and includes various fields relevant for the particular packet being sent. For instance, header 302a may include a buffer number indicating which buffer in the STI macro is to receive the packet, and a header length count indicating the length of the header.

Packet 300b also includes, for example, a header 302b and a longitudinal redundancy check 304b, as well as a payload 306 and a further longitudinal redundancy check 308. In this case, header 302b includes a payload length count indicating the length of the payload (e.g., data), in addition to the buffer number and header length count. Payload 306 is, for example, variable in length in increments of 4 bytes, and has a maximum length of 128 bytes (the memory line size).

After a packet is transmitted or when there are no packets or control information to be transmitted, an idle sequence is transmitted.

To send and receive packets over STI, each STI macro 200 includes, for example, four transmit header buffers, four receive header buffers, two transmit

payload buffers and two receive payload buffers.  
Each payload buffer is tied to one of the header  
buffers. In one embodiment, each of the header  
buffers has space for a 16 byte header, and each of  
5 the payload buffers has space for 128 bytes (a memory  
line) of payload information. The payload buffers  
are sufficient to realize the full data rate of the  
link.

10 The low level flow control provided by the STI  
macros regulate the flow of packets as they travel  
through the payload and header buffers. Some of the  
bits in the packet headers are under exclusive  
control of the STI macros and others may be used by  
the STI user layer logic (such as channel adapters  
15 118, 120, 122 and switch logic 202) for any purpose  
required.

Switch logic 202 (FIG. 2) is coupled to the STI  
macros. It is responsible for forwarding information  
received from the macros to memory and for passing  
20 information from memory to the STI macros.

When a user layer, such as the switch logic,  
wants to send a packet, it queries the STI macro to  
determine the type of buffers available. If the  
desired buffer (header and payload or only a header)  
25 is available, the user layer places its information  
on the buffer(s), and tells the STI macro to proceed.  
The STI macro sends the packet to the other end of  
STI where it is placed in the receive buffer(s). The  
STI macro at the other end informs its user layer of  
30 the arrival of the packet, and the user layer unloads  
it. After the user layer unloads the packet, the STI

macro sends a control word (e.g., an acknowledgement) back to the side of the STI originating the packet. One example of a control word is depicted in FIG. 4.

As shown, a control word 400 can be inserted  
5 anywhere either in an idle sequence 402 or within a packet 404. At the most basic level, STI macros move packets between user layers, and for each packet moved, there is a control word (acknowledgement) sent in the opposite direction after the packet has been  
10 successfully processed.

Negative acknowledgements in the form of control words, and additional control bits in the packet headers detect missing and/or damaged packets. In some cases, the STI macros can automatically  
15 retransmit a packet. This retransmission may cause the packets to be presented to the receiving user layer in an order different from the order that the sending user layer transmitted the packets. If the user layer is sensitive to the ordering, it is up to  
20 the user layer to determine the order.

Returning to FIG. 1, channel adapters 118, 120, 122 use STI as an interface to memory. Their basic operation is sending memory requests. Each memory request receives a response. For a memory fetch, the  
25 memory address is in the request and the response contains the data. For a memory store, the memory address and the data to be stored are in the request, and the response simply contains an indication of the successfulness of the store operation.

One example of a STI header for channel adapter memory requests is depicted in FIG. 5. Header 500 includes, for instance, one or more STI macro controls and routing information 502, an operation  
5 code (store or fetch) 504 and a memory address 506. The op code and memory address comprise a command.

The STI macro controls and routing information include, for example, a STI macro buffer number, a header length count, and a data length count for the  
10 memory operation.

One example of a STI header for channel adapter memory responses is depicted in FIG. 6. Header 600, which is similar to header 500, includes, for  
instance, one or more STI macro controls and routing  
15 information 602, and a memory response code 604. One response code is used to indicate a successful completion of a memory operation, while one or more other codes are used to indicate failures, such as an invalid memory address.

One example of the flow of packets for an I/O input (fetch or read) operation is depicted in FIG.  
7. For clarity purposes, low level STI acknowledgements are not shown. The flow of  
information is between a memory bus adapter 700 and a  
25 channel adapter 702. The channel adapter receives data from the peripherals (not shown in this figure) and typically, buffers this information in the adapter. Memory addresses and control information within the channel adapter is used to generate the  
30 STI headers (e.g., header 500), and the buffered data are the payloads. A packet 704, including a header

FOI b6 b7C b7D



and a payload, is sent from adapter 702 to memory bus adapter 700. A further packet 706 is also sent from the channel adapter to the memory bus adapter. In this particular example, two storage requests are  
5 outstanding at the same time. A response packet 708 (e.g., header 600) is then sent from memory bus adapter 700 to channel adapter 702. This response is paired to the request having the same STI macro buffer number as the response. Thereafter, another  
10 request packet 710 is sent from the channel adapter to the memory bus adapter. Multiple exchanges are typically required, since the peripheral data is usually more than 128 bytes. Next, two other response packets 712, 714 are sent from the memory  
15 bus adapter to the channel adapter.

One example of the flow of packets for an I/O output (store or write) operation is depicted in FIG. 8. For clarity purposes, low level STI acknowledgements are not shown. Again, the flow of  
20 information is between a memory bus adapter 800 and a channel adapter 802. With a write operation, the channel adapter is instructed to send data to the peripherals. Memory addresses and control information within the channel adapters are, once  
25 again, used to generate the STI headers. Request packets 804, 806 (e.g., header 500) are sent from adapter 802 to memory bus adapter 800 requesting data that is to be written to the peripherals. Then, a response packet 808 (e.g., header 600) is sent from  
30 the memory bus adapter to the channel adapter, which provides the data. Thereafter, another request packet 810 is sent from the channel adapter to the memory bus adapter. In response to the request

093448 06304

packets, response packets 812, 814 are sent from the memory bus adapter to the channel adapter.

Described in detail above is one embodiment of how communication between memory and peripherals within a computing system occurs. Now, the discussion turns to communication between computing systems that are coupled to one another. In particular, communication between central processing complex 102 and coupling facility 104 is described below.

In one example, communication between central processing complex 102 and coupling facility 104 is via an intersystem channel (ISC) 132 (FIG. 1). Intersystem channel 132 couples an intersystem channel adapter 122 on system 102 with an intersystem channel adapter 122 on system 104. ISC supports connections up to approximately 20 kilometers, and using fiber optic conductors, the propagation time on the link (at 20 kilometers) is approximately 100 microseconds in each direction. To minimize message latency, a minimum number of round trip acknowledgements is used. Thus, large data buffers are used for flow control.

Information is sent across an intersystem channel in frames, one example of which is depicted in FIG. 9a. A frame 900 includes a header 902, a cyclical redundancy check 904, a payload 906 and a payload cyclic redundancy check 908. Some frames do not include the payload or payload cyclic redundancy check 908. These frames are used to control the

link. Frames that include the payload are for the message commands, data and responses.

Header 902 includes, for instance, a "N" field 910 (FIG. 9b) specifying a buffer set number  
5 corresponding to a buffer set that can include, for example, a MRB, a MCB and/or data; a "R" field 912 differentiating requests (outbound) from responses; and a "D" field 914 specifying a data frame, as opposed to a command or response frame. In command  
10 and response frames, an "A" bit 916 of the header indicates that there is at least one data area associated with the message, and in data frames, it indicates that more data areas are to follow. The header also includes a "S" field 918 indicating that  
15 the data frame is the first of one or more data frames for the associated data area; and a block count 920 indicating the number of 256 byte blocks in the associated data area. In one example, commands and responses are limited to one block.

20 One example of the flow of a message from a processor running OS/390 to a processor running coupling facility code is depicted in FIG. 10a. At a minimum, a command 1000 (referred to as a message command block (MCB)) is sent from the OS/390  
25 processor to the coupling facility, and a response 1002 (referred to as a message response block (MRB)) is sent from the coupling facility back to the OS/390 processor. "A" bit 916 (FIG. 9b) is off in both  
30 either the command or the response.

TOP SECRET

The above exchange only requires a single round trip over the link, since each end of the ISC link has data buffers large enough to store the entire command and response. Assuming commands and responses that are 256 bytes and a link data transfer rate of 100 megabytes per second, the data transfer on the link takes about 5 microseconds, while the link propagation time is approximately 200 microseconds.

10        FIG. 10b depicts the flow of information between  
an OS/390 processor and the coupling facility when a  
write operation is requested. Initially, a message  
command block 1004 is sent from OS/390 to the  
coupling facility. Since data is to follow, the "A"  
15    bit is on. Thereafter, a write data frame 1006 is  
sent from the OS/390 processor to the coupling  
facility. In this example, the "A" bit is off, since  
no other data is to follow. In response, the  
coupling facility sends a message response block 1008  
20    to the OS/390 processor.

In the above example, the write data frame immediately follows the MCB; no additional exchanges are required on the link before the data frame is sent. This is due to the large buffers of ISC that  
25    are capable of temporarily storing the frames until  
the memory addresses of the payload can be  
determined.

FIG. 10c depicts the flow of information between an OS/390 processor and the coupling facility when a  
30    read operation is requested. A message command block  
1010 is sent from the OS/390 processor to the

098443 0504  
10590 8842860

coupling facility. Since no data is to follow (this is a read request), "A" bit 916 is off. In response to the command, a read data frame 1012 is sent from the coupling facility to the OS/390 processor.

5 Further, a message response block 1014 is also sent from the coupling facility. In this scenario, the "A" bit is on; however, instead of indicating that data is to follow, it is used by the link to detect whether any data frames have been lost.

10 The added time for data transfer, assuming 4096 bytes of data, is about 40 microseconds. Additional 4096 byte data buffers at each end of the link maintain the single round trip. As more data (greater than 4096 bytes) is required as part of the  
15 message, more buffering is required at each end of the ISC link to maintain the single round trip flow control. The expense of large buffers eventually becomes prohibitive, and the flow control is modified to add intermediate acknowledgement exchanges to  
20 throttle the data transfer preventing the buffers from overrunning. These extra exchanges, however, significantly increase message latency.

One example of the flow control for a write operation using intermediate acknowledgements is  
25 depicted in FIG. 11a. Initially, a message control block 1100 is sent from the OS/390 processor to the coupling facility. In this example, the "A" bit is on indicating that data is to follow. Next, a data frame 1102 is sent to the coupling facility. The  
30 data frame also has the "A" bit on instructing the other end of the link to send a link acknowledgement 1104 when the data buffer area is ready to receive

PO9-98-109-21

the next data area. After link acknowledgement 1104 is received by the OS/390 processor, at least one other data frame is sent (e.g., data frames 1106, 1108). Since data frame 1108 has the "A" bit off, the coupling facility knows that it is the last data frame and instead of sending an acknowledgement, it sends a message response block 1110.

One example of the flow control for a read operation using intermediate acknowledgements is depicted in FIG. 11b. Initially, a message control block 1112 is sent from the OS/390 processor to the coupling facility. The MCB has the "A" bit off. Thereafter, the coupling facility sends back a data frame 1114. This data frame has the "A" bit on indicating that an acknowledgement is to be sent from the OS/390 processor to the coupling facility. Thus, the OS/390 processor sends an acknowledgement 1116. When the coupling facility receives this acknowledgement, it knows that it can send more data. Thus, data frame 1118 is sent. Then, when the coupling facility receives the acknowledgement for that data frame, another data frame 1120 is sent to the OS/390 processor. This data frame has the "A" bit turned off, so the next frame sent is a message response block frame 1122 from the coupling facility to the OS/390 processor.

In addition to or in lieu of the above ISC communication path between central processing complex 102 and coupling facility 104, the present invention provides another, more direct, communication path between the central processing complex and the coupling facility. This communication path is

referred to as an integrated cluster bus (ICB) 134  
(FIG. 1).

Integrated cluster bus 134 couples the memory  
bus adapter of the central processing complex with  
5 the memory bus adapter of the coupling facility over  
a self-timed interface link. In order for the self-  
timed interface to be adapted to be used as an  
integrated cluster bus (i.e., in order for the self-  
timed interface to be used as a channel), the channel  
10 functions of a direct memory adapter engine and  
memory protection are added to the memory bus  
adapter, in accordance with the principles of the  
present invention. Further, the STI packet  
definition and the STI flow control are expanded in  
15 order to accommodate the integrated cluster bus (or  
channel functions), as described herein.

The addition of control logic to the memory bus  
adapter 114 to support integrated cluster bus flow  
control is depicted in FIG. 12. As with the memory  
20 bus adapter of FIG. 2, the memory bus adapter of FIG.  
12 includes STI macros 1200 and switch logic 1202,  
which are similar to macros 200 and logic 202,  
described above. However, memory bus adapter 114 of  
FIG. 12 also includes ICB control 1204. In this  
25 embodiment, there is an ICB control 1204 coupled to  
each STI macro 1200. However, for any STI link that  
is coupled to a channel adapter, such as adapter 118,  
120, 122, the ICB control is disabled and bypassed,  
as shown by the arrows in FIG. 12. Thus, since only  
30 one STI link is coupled between memory bus adapters,  
only one of the ICB controls is enabled (no arrows  
going through the last ICB control).

Integrated cluster bus control 1204 is the implementation of the ICB link that couples the memory bus adapters. It includes, for instance, part of intersystem channel adapter 122 direct memory address engine (in order to emulate ISC channel adapter 122) and circuitry for a self-timed interface user layer protocol to support S/390 message passing, in accordance with the principles of the present invention. ICB control 1204 has both main attributes of a S/390 channel. It is a DMA engine and it protects memory from being arbitrarily overwritten by the attached peripherals. (An example of a DMA engine is described in co-pending, commonly assigned U.S. Patent Application No. 08/788,755, Gregg et al. entitled "Speeding-Up Communication Rates on Links Transferring Data Structures By a Method of Handling Scatter/Gather of Storage Blocks In Commanded Computer Systems," filed on January 24, 1997, which is hereby incorporated herein by reference in its entirety.)

One example of a header used for ICB is depicted in FIG. 13. In one example, a header 1300 (which has a length of, for instance, 8 bytes) includes an "I" bit 1302 identifying the packet as an ICB packet, as opposed to other control packets used to set interrupt bits for such functions as canceling or terminating messages. If the "I" field indicates an ICB packet, then a "N" field 1304, a "R" field 1306 and a "D" field 1308 of the header are used to identify the buffer set and appropriate areas within the buffer set. In particular, "N" field 1304 specifies a buffer set number; "R" field 1306 differentiates requests (outbound) from responses;



and D field 1308 specifies a data frame, as opposed to a command or response frame.

In addition to the above fields, in accordance with the principles of the present invention, ICB header 1300 also includes a DR field 1310 identifying a Data Request packet for the buffer set specified by "N" field 1304; a "C" field 1312 (continue field) indicating that more packets are to be transmitted to complete the data for the buffer area specified by "N" field 1304, "R" field 1306 and "D" field 1308; and a SEQ field 1314 providing a sequence count used to put the received packets in the correct order. The DR, C, and SEQ fields are described in further detail below.

Additional details of ICB control 1204 are described with reference to FIG. 14. In order to take advantage of ICB, the ICB control is enabled at both ends of the appropriate STI link (i.e., at the OS/390 side of the link and at the coupling facility side). When ICB is enabled, "I" field 1302 (FIG. 13) of each packet header received from the STI macro is examined by circuit 1400 of ICB control 1204. Circuit 1400 determines if the "I" field is equal to ICB.

If the "I" field is not equal to ICB, then the packet is not an ICB packet, so the packet is simply passed up to switch logic 1202 via an OR circuit 1402, where it is treated like the packets generated by channel adapters.

However, should the "I" field be equal to ICB, then the packet is an ICB packet and the header is sent to a packet header handler 1404. This circuitry includes the DMA engine, flow control logic and the memory request/response header modification logic of the present invention. This logic determines, for example, whether the ICB packet is for a MCB, MRB or data area. If it is for a MCB, MRB or data area, as indicated by "R" field 1306 and "D" field 1308, the state of the buffer area is checked to see if it is ready to accept data. Should the buffer area be ready, the ICB header is expanded to, for instance, 16 bytes, and is modified into a memory store command STI header (e.g., header 500 (FIG. 5)). The memory address is generated from address registers and control and state information in the DMA engine.

After the new header is generated, the payload is taken from the STI macro and appended to the new header by OR circuit 1402 (also referred to as a combine circuit). The ICB packet header is thus transformed into a STI memory store packet header. Since the packet sent to the MBA switch via OR circuit 1402 is a normal memory store operation, there is a response packet sent back towards the STI macro. These packets are intercepted by an ICB? circuit 1406 and sent to packet header handler 1404. These packets are not sent to the STI macro. If the memory response reason code indicates a successful operation, the DMA engine marks the memory storage operation as finished, allowing it to send another memory store operation when the next ICB packet is received. More than one memory store operation may be in process concurrently.

09874438 "050501  
105090" 88442860

On the other hand, if the packet received from the switch is not an ICB packet, INQUIRY 1406, then the packet is passed on to the STI macro via OR circuit 1408.

5           In the outbound direction, when the packet header handler wants to send an ICB packet over STI, it first generates a memory fetch packet using memory addresses from the DMA engine. Then, it sends the memory request to the MBA switch through OR circuit  
10 1402. When the memory response packet is detected by ICB? circuit 1406, the header is sent to the packet header handler where the "N", "R", "D", DR, "C" and SEQ fields are added to the first 8 bytes of the header. The second 8 bytes contain the memory  
15 request response code, and if this code indicates a successful memory request, the packet header handler sends the new ICB header to OR circuit 1408 followed by the payload from circuit 1406. The packet is then forwarded to the STI macro via OR circuit 1408.

20           In addition to the above, a connection 1410 exists between the processors and ICB control 1204, in order for the processors and ICB control to communicate more directly. This is shown in the following examples in which messages are sent via the  
25 ICB control.

FIG. 15 depicts a write data example in which a message is sent from an OS/390 processor 1500 to a coupling facility processor 1502. As described, the message proceeds from processor 1500 through ICB  
30 control 1504 over an ICB STI link 1506 and through ICB control 1508 to coupling facility processor 1502.

05037443 060504

For clarity purposes, only the ICB packets are shown in this example, as well as the example depicted in FIG. 16. The low level STI link acknowledgements are not shown.

5           Initially, OS/390 processor 1500 sends a message by issuing a Send Message (SMSG) instruction 1510, which causes a start command to be sent to ICB control 1504, STEP 1512. This command instructs the ICB control to fetch the MCB and send it over ICB STI  
10 1506 in one or more packets, and to begin preparing the MRB, STEP 1513. (One example of preparing the MRB is described in detail in co-pending, commonly assigned U.S. Patent Application No. 08/864,583, Gregg et al., entitled "Intersystem Advanced Mode  
15 Sending Method and Mechanism", filed May 28, 1997, which is hereby incorporated herein by reference in its entirety.)

          In this example, two MCB packets 1514a, 1514b are used, each having a maximum of 128 bytes of data.  
20 The "C" bit in MCB packet 1514a indicates to the control logic that another packet is forthcoming. ICB control 1508 receives the packets via ICB STI link 1506 and stores them, STEP 1520. It also indicates to the coupling facility that work is to be  
25 performed, STEP 1522.

          After the MCB packets have been received by coupling facility processor 1502, the coupling facility code detects the command and determines whether or not there is write data associated with  
30 the command, STEP 1524. Should there be write data, as in this case, the data packets are not sent to the

09074438-060501

coupling facility until the coupling facility code is ready to receive them (e.g., has determined the memory addresses). This is because the STI links do not have the expense and complexity of large buffers.

5 Without the large buffers, the data packets cannot be sent until the coupling facility is ready.

When the coupling facility code receives the MCB packets and determines there is write data, the coupling facility code determines the memory

10 addresses for the data and issues a Move Channel Buffer Data Multiple (MCBDM) instruction to move the data, STEP 1526. The code that performs the MCBDM instruction sends a read list command 1528 to ICB control 1508. (The fetching of addresses for data

15 areas is further described in co-pending, commonly assigned U.S. Patent Application No. 08/788,755, Gregg et al. entitled "Speeding-Up Communication Rates on Links Transferring Data Structures By a Method of Handling Scatter/Gather of Storage Blocks

20 In Commanded Computer Systems," filed on January 24, 1997, which is hereby incorporated herein by reference in its entirety.) The ICB control then fetches the first address in the list, STEP 1530, and sends a Data Request over STI 1506 to ICB control

25 1504, STEP 1532.

In addition to the above, when the SMSG instruction is issued, the code on the OS/390 processor sends a write list command to ICB control 1504, if data transfer is required as part of the

30 message, STEP 1534. If no data is associated with the message, the list command is not issued to the ICB control and, therefore, no data packets are sent

05874438-060504

over STI. However, as a result of the write list command, ICB control 1504 fetches the first entry in the list of addresses pointing to the data to be moved, STEP 1536.

5           After ICB control 1504 has received the write list command from processor 1500 and has fetched the first address in the list, and after ICB control 1504 has received the Data Request packet from the other end of the STI, ICB control 1504 can proceed to  
10           transmit data packets. (The Data Request packet can be received before or after ICB control 1504 has fetched the first address of the list.) In particular, ICB control 1504 fetches data from the memory addresses that are provided, STEP 1538, and  
15           places the data in one or more data packets, which are transmitted from ICB control 1504 to ICB control 1508, STEP 1539. When ICB control 1508 receives the packets, it stores them, STEP 1540, and fetches the next list entry, STEP 1542. The ICB control knows  
20           there is more data to be fetched, since the "C" bit is on in the last data packet received. Processing then continues, as described above.

          When there is no more data to be stored (the "C" bit of the last data packet is off indicating the end  
25           of a buffer area), ICB control 1508 indicates completion, STEP 1550, which is detected by the coupling facility, STEP 1552. Thus, the coupling facility triggers ICB control 1508 to fetch an MRB, STEP 1554. The ICB control fetches the MRB, STEP  
30           1556, and transmits one or more MRB packets to ICB control 1504. ICB control 1504 receives the packets and stores them, STEP 1558. When ICB control 1504

PO9-98-109-30

detects that it has received all of the MRB packets,  
it sends a done indication to the OS/390 processor,  
STEP 1560.

FIG. 16 depicts a read data example in which a  
5 message is sent from an OS/390 processor 1600 to a  
coupling facility processor 1602. As described, the  
message proceeds from processor 1600 through ICB  
control 1604 over an ICB STI link 1606 and through  
ICB control 1608 to coupling facility processor 1602.  
10 Again, for clarity purposes, only the ICB packets are  
shown in this example.

Initially, OS/390 processor 1600 sends a message  
by issuing a Send Message (SMSG) instruction 1610,  
which causes a start command to be sent to ICB  
15 control 1604, STEP 1612. This command instructs the  
ICB control to fetch the MCB and send it over ICB STI  
1606 in one or more packets, and to begin preparing  
the MRB, STEP 1613. In this example, two MCB packets  
1614a, 1614b are used, each having a maximum of 128  
20 bytes of data. ICB control 1608 receives the packets  
and stores them, STEP 1620. It also indicates to the  
coupling facility that work is to be performed, STEP  
1622.

After the MCB packets have been received by  
25 coupling facility processor 1602, the coupling  
facility code detects the command, STEP 1624,  
determines the memory addresses for the data and  
issues a Move Channel Buffer Data Multiple (MCBDM)  
instruction to move the data, STEP 1626. The code  
30 that performs the MCBDM instruction sends a write  
list command to ICB control 1608, STEP 1628. The ICB

PO9-98-109-31-16

control then fetches the first address in the list,  
STEP 1630.

In addition to the above, when the SMSG  
instruction is issued, the code on the OS/390  
5 processor sends a read list command to ICB control  
1604, STEP 1636. As a result of the read list  
command, ICB control 1604 fetches the first entry in  
the list of addresses pointing to the data to be  
moved, STEP 1638. Further, it sends a Data Request  
10 packet to the other end of the STI, STEP 1640.

After ICB control 1608 has fetched the first  
address in the list, and after it has received the  
Data Request packet from the other end of the STI,  
ICB control 1608 can proceed to transmit data  
15 packets. (The Data Request packet can be received  
before or after ICB control 1608 has fetched the  
first address of the list.) In particular, ICB  
control 1608 fetches data from the memory addresses  
that are provided, STEP 1642, and places the data in  
20 one or more data packets, which are transmitted from  
ICB control 1608 to ICB control 1604, STEP 1644.  
When ICB control 1604 receives the packets, it stores  
them, STEP 1646, and fetches the next list entry,  
STEP 1648. Processing then continues, as described  
25 above.

When there is no more data to be fetched, ICB  
control 1608 indicates completion, STEP 1650, which  
is detected by the coupling facility, STEP 1652.  
Thus, the coupling facility triggers ICB control 1608  
30 to fetch an MRB, STEP 1654. The ICB control fetches  
the MRB, STEP 1656, and transmits one or more MRB

105090" 88442860



packets to ICB control 1604. ICB control 1604 receives the packets and stores them, STEP 1658. When ICB control 1604 detects that it has received all of the MRB packets, it sends a done indication to  
5 the OS/390 processor, STEP 1660.

In the examples described above, the ICB control logic is responsible for performing a number of functions, including, for instance, executing list commands and interpreting Data Request and data  
10 packets. Additionally, the ICB control logic is responsible for determining various error conditions. Further details of how these functions are performed by the ICB control logic on both sides of the link (e.g., sender and receiver) are described in detail  
15 with reference to FIG. 17.

In one example, initially, a command is received by the ICB control from an OS/390 processor via connection 1410 (FIG. 14), STEP 1700. This command is checked to determine whether it is a read list  
20 command, INQUIRY 1702. When a list command for a read operation is received, a LIST READ LATCH located in the ICB control is tested to see if a list command for a read operation has already been received, INQUIRY 1704. If a read list command has already  
25 been received, then an error, referred to as a LIST READ DOUBLE ERROR, is set, STEP 1706 and processing is complete, STEP 1708.

If, however, the LIST READ LATCH indicates that a read list command has not previously been received,  
30 then the LIST READ LATCH is set (e.g., set to one), STEP 1710, and a further determination is made as to

09874438-060504

whether a Data Request packet has previously been received, INQUIRY 1712. In one example, this determination is made by checking a D\_REQ RCVD LATCH located in the ICB control. Since this is a read  
5 operation, a Data Request packet should not have been received. Thus, if the latch indicates one has been received, than a D\_REQ ERROR is set to one, STEP 1714, and processing is complete, STEP 1708.

10 If a Data Request packet has not been received, the ICB control logic sends a Data Request packet to the other end of the link, STEP 1716, and processing is once again complete, STEP 1708.

Returning to INQUIRY 1702, if the list command is not for a read operation, then a determination is  
15 made as to whether the command is for a write operation, INQUIRY 1718. If the list command is not for a read or write operation, then processing completes, STEP 1720. On the other hand, if the list command is for a write operation, then a LIST WRITE  
20 LATCH is tested to determine whether a list command for a write operation has already been received, INQUIRY 1722.

Should a write list command have previously been received, then a LIST WRITE DOUBLE ERROR is set  
25 (e.g., to one), STEP 1724, and processing is complete, STEP 1708. However, if a list command for a write operation has not been received, then the LIST WRITE LATCH is set to one, STEP 1726, and the D\_REQ RCVD LATCH is tested, INQUIRY 1728.

09074498 060501

If a Data Request packet is received, then processing continues by sending the data packets to the other end of the link, STEP 1730. However, if a Data Request packet has not been received, then  
5 processing ends, STEP 1708.

When an ICB packet is received by the ICB control logic, STEP 1800 (FIG. 18), it is tested to determine if it is a Data Request packet, INQUIRY 1802. If the "DR" field of the header indicates that  
10 the packet is a Data Request packet, then the D\_REQ RCVD LATCH is tested to determine if a Data Request packet was previously received, INQUIRY 1804. Should the latch be set, then D\_REQ DOUBLE ERROR is set to one, STEP 1806, and processing is complete, STEP  
15 1808.

On the other hand, if the D\_REQ RCVD LATCH is not set, then the latch is set, STEP 1810, and a LIST READ LATCH is then tested, INQUIRY 1812. If the LIST READ LATCH is set indicating that this is a read  
20 operation, then a D\_REQ ERROR is set, since a Data Request packet should not have been received, STEP 1814. Processing is then complete, STEP 1808.

If, however, the list read latch is not set, then a determination is made as to whether a write  
25 operation is in progress, INQUIRY 1816. This determination is made by checking the LIST WRITE LATCH. If it is set, then a write operation is in progress. When a write operation is in progress, the ICB control logic starts sending data packets, STEP  
30 1818. When the LIST WRITE LATCH is not set, then processing ends.

105090" 88442860

Returning to INQUIRY 1802, if the received packet is not a Data Request packet, then a determination is made as to whether the packet is a data packet, INQUIRY 1822. If the packet is not a data packet, then processing ends. However, if the packet is a data packet, then the LIST READ LATCH is checked to determine if a read operation is in progress, INQUIRY 1824. When a read operation is not in progress, a DATA ERROR is set, since data is being received prior to sending the Data Request packet, STEP 1826.

However, if a read operation is in process, then the packet payload is sent to memory, STEP 1828, and processing is complete, STEP 1808.

In addition to the above-described error checking, the ICB control logic also detects protocol errors involving the "C" bit of the ICB headers. As described above, this bit is used to indicate an end of a buffer area. That is, for MCB, MRB, or data areas, the "C" bit is used to determine the length of those areas. If the bit is on, more data is to follow. If it is off, there is no more data and the end of a buffer area is indicated. (In this embodiment, no other fields are necessary to indicate the length of a buffer area, since the low level STI guarantees delivery of all packets.)

Based on the foregoing, once a packet is received with its "C" bit off, the MCB, MRB, or data buffer area state is set to full. If another packet is received for a full buffer area, a C bit protocol

error is set. The full state is reset at the beginning of the next message.

In accordance with the principles of the present invention, ICB data transfer does not pair each sent packet with a response packet. Thus, the bandwidth required for response packets in other data transfer modes (e.g., channel adapters) is not needed. Therefore, this bandwidth is used to send more payload information. As a result, the ICB packets are streamed together, and the arrival order is important in determining where in memory each packet payload is stored.

The in-order packet arrival of STI is disturbed by STI's automatic packet retry. Through the use of negative acknowledgements, the STI macros resend damaged or lost packets, and because multiple packets and their link acknowledgements can exist on the link at the same time, the apparent order of the arrived packets to the ICB control may not be in the same order that the packets were transmitted.

Thus, in accordance with the principles of the present invention, the STI headers for ICB include a SEQ field 1314 (FIG. 13) to determine the order of the STI packets. When the ICB control sends a packet to the STI macro, it places the contents of its sequence counter in the SEQ field of the header and increments the counter. Each time a packet is received, the ICB control compares the SEQ field to its sequence counter. If the counts are the same, the packet is moved from the STI macro to memory at the next address. If the counts are not the same

(because a packet was resent), the ICB control leaves the packet in the STI macro and waits for the next packet to be received. The payload of the next packet is then stored in memory, and the ICB control compares the sequence counts of the previously received packet. It expects, in this embodiment, that this time it will be the same, since no more than two packets are sent at any one time. The packet is then stored in memory, and the sequence counter is again incremented.

Further details of the sequence number handling is described below with reference to FIG. 19. When a packet is received off the link by the STI macro, STEP 1900, SEQ field 1314 in the header (referred to in FIG. 19 as HDRSEQ) is compared to the receiver's sequence count (RCVSEQ), INQUIRY 1902. If the counts do not match, then the control logic determines whether the header sequence count is one more than the receiver's sequence count, INQUIRY 1904.

When the header sequence count is not one more than the receiver's count, a SEQ ERROR is set, STEP 1906. However, if the header sequence count is one more than the receiver's sequence count, then the packet is left in the STI macro packet buffer, and the ICB control sets a NEXT AVAILABLE PACKET indicator to one indicating a packet is waiting to be processed, STEP 1908.

Returning to INQUIRY 1902, if the counts match, the packet is processed, STEP 1910. For example, the payload is transmitted to memory. Additionally, the

receiver's sequence count is incremented by one, STEP 1912.

5       Thereafter, the ICB control logic checks the  
NEXT PACKET AVAILABLE indicator to determine if a  
previously out of sequence packet is in the STI macro  
packet buffer, INQUIRY 1914. If there is no out of  
sequence packet, then processing ends, STEP 1915.  
However, if a packet is out of sequence, then that  
packet is processed by sending the payload to memory,  
10   STEP 1916. Additionally, RCVSEQ is incremented by  
one, STEP 1918, and NEXT PACKET AVAILABLE is reset to  
zero, STEP 1920.

15       In an alternate embodiment, the SEQ field in the  
header can be used to modify or calculate the memory  
address to store the payload of the packet for a MCB,  
MRB or data area. In this embodiment, the ICB  
control keeps track of which payloads are stored to  
ensure that there are no holes in the data area in  
memory. If the packet is a Data Request packet, the  
20   packet can be processed immediately, since ordering  
of this packet with respect to other ICB packets is  
not necessary.

25       While the ICB control sequence numbers keep ICB  
packets in the correct order, the other control  
packets (indicated by "I" bit 1302 (FIG. 13) in the  
header set to zero) that bypass the ICB control logic  
are to be ordered with respect to the ICB packets.  
For example, if a message is to be canceled,  
terminated or withdrawn by the sender, the control  
30   packet initiating the cancel operation cannot get  
ahead of the message. As shown above, the resending

105090" 88442860

of damaged packets may allow the control packet to get ahead of the message. The sequence numbers in the ICB message packets are not used by the control packets. A quiesce function in the ICB control logic  
5 is provided to keep control packets in the correct order. Before the code issues a control packet, it commands the ICB control logic to quiesce and restart the link. When the quiesce command is received by the ICB control logic, it quits sending packets and  
10 waits for all outstanding packets to be acknowledged at the STI link level. Once the ICB packets are all acknowledged, the ICB control logic knows that there are no damaged or lost packets being resent by the STI macros. As soon as the ICB control logic detects  
15 that all packets have been acknowledged, it resumes sending more packets. After the code detects that the quiesce operation has completed, it knows that it can send a control packet. The quiesce function provides a serialization point for the cancel  
20 operation.

The ICB control logic has additional state machines for checking the order of commands received from processors, which is described with reference to FIG. 20.

25 Initially, the ICB control logic receives a command from a processor, STEP 2000, and that command is tested to determine whether it is a start command, INQUIRY 2002. If it is a start command, a START LATCH is tested to determine whether a start command  
30 has already been received, INQUIRY 2004. When a start command has already been received, then a START

0987448-0650-0650



DOUBLE ERROR is set, STEP 2006, and processing is complete, STEP 2008.

However, if a start command has not been previously sent, then the START LATCH is set, STEP 2010, and the MCB is sent to the other end of the link, STEP 2012.

Returning to INQUIRY 2002, if the command from the processor is not a start command, then a determination is made as to whether the command is a list command (read or write operation), INQUIRY 2014. If the command is not a list command, then processing ends, STEP 2015. However, if the command is a list command, the START LATCH is tested to make sure it is on before the list command is received, INQUIRY 2016. When the START LATCH is on, processing ends, STEP 2015. On the other hand, when it is not on, a LIST ERROR is set, STEP 2018, and then, processing is complete, STEP 2008.

Further, the ICB control logic has additional state machines for checking the order of packets received from the other end of the link, which is described with reference to FIG. 21.

When a packet is received by the ICB control logic, STEP 2100, a determination is made as to whether the packet is a MRB, INQUIRY 2102. If the received packet is not a MRB, processing ends, STEP 2103. However, if the received packet is a MRB, the START LATCH is tested, INQUIRY 2104. Should the latch be off, then MRB ERROR is set, STEP 2106. However, if it is on, the MRB packet payload is sent

105090" 88442860

to memory, STEP 2108, and the START LATCH is set to zero, STEP 2110.

It will be apparent to those skilled in the art that although a single buffer set is described above, multiple buffer sets may be transferring packets in an interleaved or multiplexed fashion. Each buffer set has separate state machines and controls that operate independently.

Described in detail above are ways to control the flow of information over a link that is being used as a channel. One technique described is the addition of a Data Request packet that lets the sender know when the receiver is ready to receive the data. Advantageously, the exchange of the Data Request packets does not significantly add to message latency, since, in the above examples, STI distances are limited to about 10 meters. Further, these packets are processed by fast hardware state machines located in the ICB control.

20           Without the Data Request packets, the throughput  
over STI would be greatly reduced. For example, in  
the write example described with reference to FIG.  
15, if ICB control 1504 started to immediately send  
data packets after it received the list command  
25   (1534) from processor 1500, the first two data  
packets sent on STI would be buffered in the STI  
macro at the coupling facility end of the STI. These  
packets could not be moved to memory, since the  
address(es) is not yet known. Thus, all other  
30   traffic for all other buffer sets would stop, waiting  
for these two data packets to be stored. Eventually,

the coupling facility code would detect the command, determine the address(es), and send the list command to ICB control 1508. Only then could the contents of the STI macro data buffers be stored in memory, freeing them for the transfer of more packets. It typically takes several tens of microseconds for the coupling facility code to detect the message command and determine the memory addresses. At STI speeds of hundreds of bytes per microsecond, many thousands of bytes of bandwidth would be wasted by each message. Without the Data Request protocol, the throughput of ICB suffers greatly.

In ICB, the link acknowledgement sequence shown in FIGs. 11a-11b for ISC is eliminated. The low level flow control is the STI macro buffer acknowledgements. Any number of data packets can be transferred without the higher level link acknowledgements required by ISC. For example in FIG. 15, the first data area is transmitted by data packets 1539, and the second data area is transmitted by data packets 1537. There is no higher level link acknowledgement between those data areas. The same situation exists for the read operation of FIG. 16, where there is no higher level link acknowledgement between data area packets 1644 and 1645.

The present invention can be included in an article of manufacture (e.g., one or more computer program products) having, for instance, computer usable media. The media has embodied therein, for instance, computer readable program code means for providing and facilitating the capabilities of the present invention. The article of manufacture can be

PO9-98-109-43

included as a part of a computer system or sold separately.

5      Additionally, at least one program storage device readable by a machine, tangibly embodying at least one program of instructions executable by the machine to perform the capabilities of the present invention can be provided.

10      The flow diagrams depicted herein are just exemplary. There may be many variations to these diagrams or the steps (or operations) described therein without departing from the spirit of the invention. For instance, the steps may be performed in a differing order, or steps may be added, deleted or modified. All of these variations are considered  
15      a part of the claimed invention.

20      Although preferred embodiments have been depicted and described in detail herein, it will be apparent to those skilled in the relevant art that various modifications, additions, substitutions and the like can be made without departing from the spirit of the invention and these are therefore considered to be within the scope of the invention as defined in the following claims.

105990 88472860